

Automating your Iris proofs with Diaframe

Ike Mulder

Radboud University Nijmegen
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Scope

Automation for *fine-grained concurrency*

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Automation for *fine-grained concurrency*:

- ▶ standard WP goals
- ▶ support for invariants $\boxed{P}^{\mathcal{N}}$
- ▶ support for ghost state \boxed{a}^{γ}

Scope

Automation that can be used *interactively*:

- ▶ no global backtracking
- ▶ extensible in language, ghost theory

Diaframe

- ▶ Plugin for Iris
- ▶ tactic based-automation: `iStepsS`

Example

```
Definition parallel_add: expr :=  
  let: "r" := ref #0 in  
  (FAA "r" #2) ||| (FAA "r" #2);;  
  !"r".
```

Prove:

$$\{\text{True}\} \text{parallel_add} \{v, \ulcorner v = 4 \urcorner\}$$

Example

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Definition parallel_add: expr :=  
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```

$\{\text{True}\}$ parallel_add $\{v, \lceil v = 4 \rceil\}$ proof:

1. execute ref
2. allocate invariant
3. use specification of |||
4. verify left & right thread
5. verify load

Example

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Definition parallel_add: expr :=  
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$\{\text{True}\}$ parallel_add $\{v, \lceil v = 4 \rceil\}$ proof:

1. ~~execute ref~~ run automation
2. allocate invariant
3. use specification of |||
4. ~~verify left & right thread~~ run automation
5. ~~verify load~~ run automation

Example

```
Definition parallel_add: expr :=  
  let: "r" := ref #0 in  
  (FAA "r" #2) ||| (FAA "r" #2);;  
  !"r".
```

{True} parallel_add {v, $\ulcorner v = 4 \urcorner$ } proof:

1. run automation
2. allocate invariant
3. use specification of |||
4. run automation
5. run automation

about $\pm 75\%$ shorter

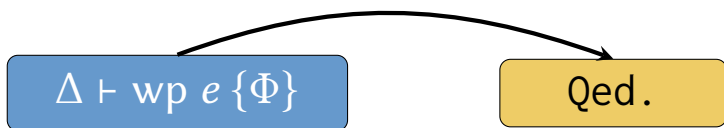
Diaframe approach

$\Delta \vdash \text{wp } e \{ \Phi \}$

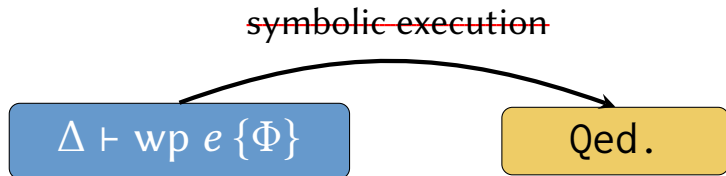
Qed.

Diaframe approach

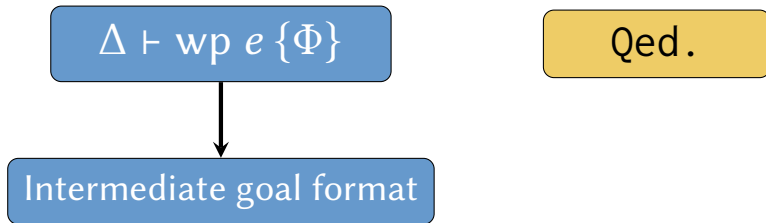
symbolic execution



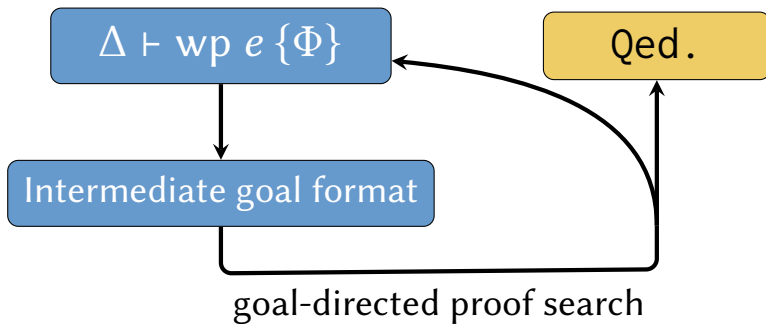
Diaframe approach



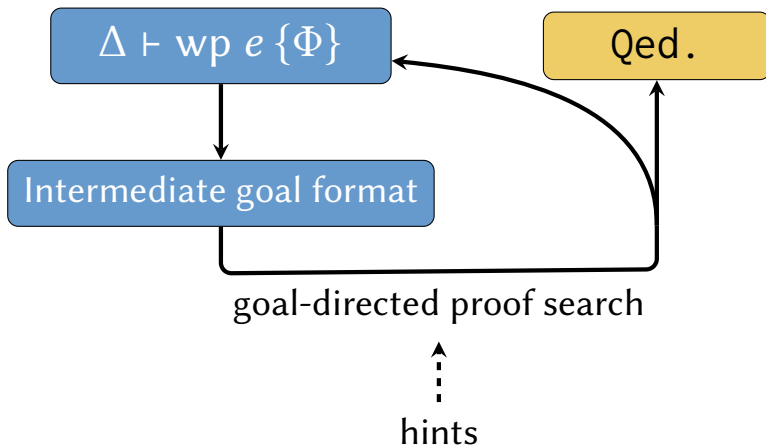
Diaframe approach



Diaframe approach



Diaframe approach

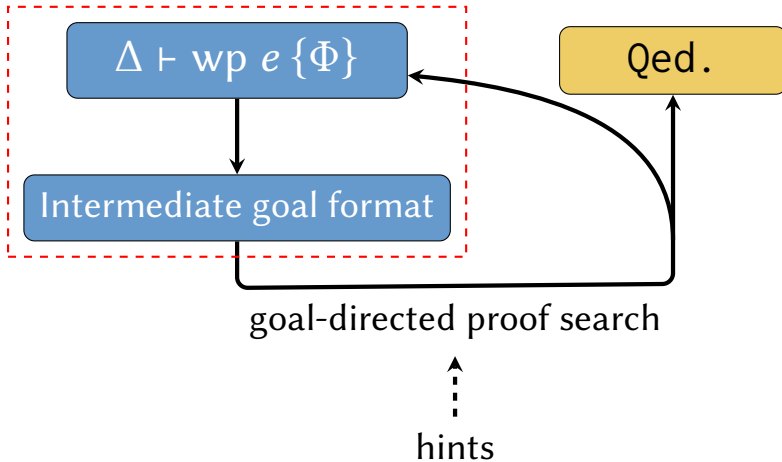


Diaframe: results

Verified 24 examples from the literature

Comparable proof burden to automated tools, but *foundational*

Caper, Voila, Starling



Goal format

$$\ell \mapsto v \vdash \text{wp } !\ell \{ \Phi \}$$

$$\boxed{\exists v. \ell \mapsto v}^{\mathcal{N}} \vdash \text{wp } !\ell \{ \Phi \}$$

wp_load ✓

wp_load ?

Goal format

$$\ell \mapsto v \vdash \text{wp } !\ell \{ \Phi \}$$

wp_load ✓

$$\boxed{\exists v. \ell \mapsto v}^{\mathcal{N}} \vdash \text{wp } !\ell \{ \Phi \}$$

wp_load ✗

hard to determine directly what invariant to open

Goal format

$$\frac{\Delta \vdash \top \stackrel{?}{\Rightarrow} \exists u. \ell \mapsto u * (\ell \mapsto u * \stackrel{?}{\Rightarrow} \top \Phi u)}{\Delta \vdash \text{wp } !\ell \{ \Phi \}}$$

Goal format

can still open invariants

$$\frac{\Delta \vdash \boxed{\top \stackrel{?}{\Rightarrow}} \exists u. \ell \mapsto u * (\ell \mapsto u * \stackrel{?}{\Rightarrow} \top \Phi u)}{\Delta \vdash \text{wp } !\ell \{ \Phi \}}$$

Goal format

for when witness is inside invariant

$$\frac{\Delta \vdash \top \stackrel{?}{\Rightarrow} \exists u. \ell \mapsto u * (\ell \mapsto u * \stackrel{?}{\Rightarrow} \top \Phi u)}{\Delta \vdash \text{wp } !\ell \{ \Phi \}}$$

Goal format

find hypothesis that can make progress

$$\Delta, \ell \mapsto v \vdash \top \stackrel{?}{\Rightarrow} \exists u. \ell \mapsto u * G$$

Goal format

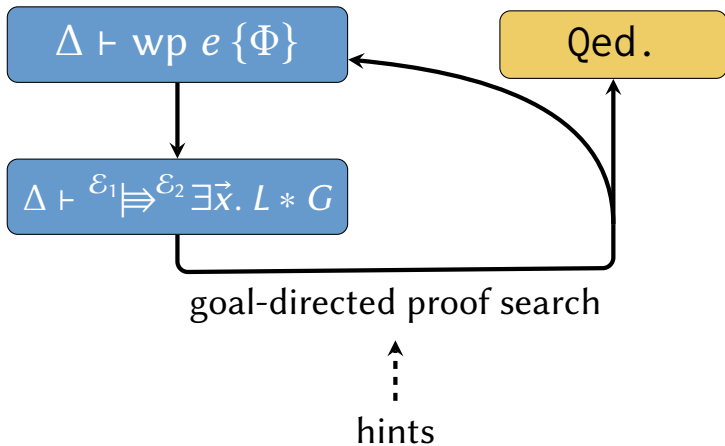
find hypothesis that can make progress

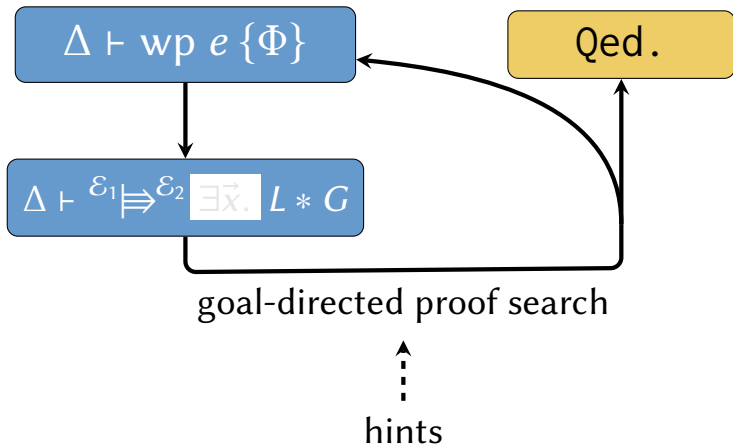
$$\Delta, \boxed{\exists v. \ell \mapsto v}^{\mathcal{N}} \vdash^{\top} \stackrel{?}{\Rightarrow} \exists u. \ell \mapsto u * G$$

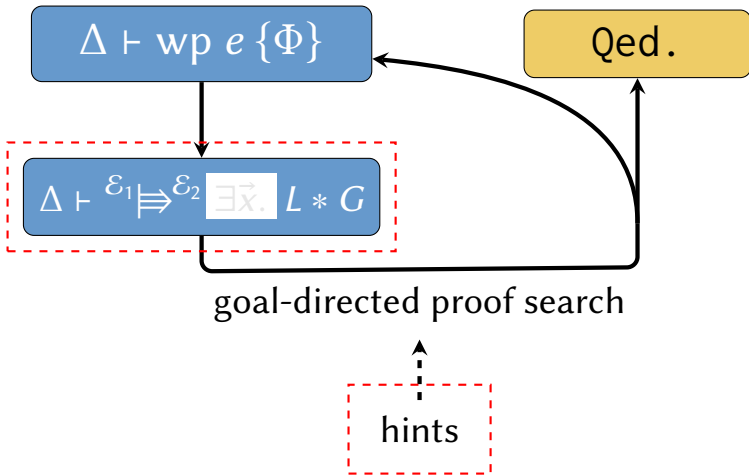
Goal format

find hypothesis that can make progress

$$\Delta, \boxed{\exists v. \ell \mapsto v * B}^{\mathcal{N}} \vdash^{\top} \stackrel{?}{\Rightarrow} \exists u. \ell \mapsto u * G$$







Hints

definition?

how to detect?

H can make progress on A ??

$$\Delta, H \vdash \varepsilon_1 \stackrel{\text{def}}{\Rightarrow} \varepsilon_2 A * G$$

Hints

A	can make progress on	A
$(B \multimap A)$	can make progress on	A
$\boxed{A * B}^{\mathcal{N}}$	can make progress on	$\triangleright A$
$\boxed{\bullet a}^{\gamma}$	can make progress on	$\boxed{\bullet a'}^{\gamma}$

Hints

$$\begin{array}{ccc} A & \vdash & A \\ (B \multimap A) * B & \vdash & A \\ \boxed{A * B}^{\mathcal{N}} & \vdash^{\mathcal{T}} \Rightarrow^{\mathcal{T} \setminus \mathcal{N}} \triangleright A * (\triangleright B * \dots) \\ \bullet a^{\gamma} & * \circ b^{\gamma} \vdash \Rightarrow & \bullet a'^{\gamma} * \circ b'^{\gamma} \end{array}$$

Hints

H can make progress on A if

$$H * L \vdash^{\mathcal{E}_1} \Rightarrow^{\mathcal{E}_2} A * U$$

for some **sidecondition** L
residue U

Hints

H can make progress on A if

$$H * [L] \Vdash [\mathcal{E}_1 \Rightarrow \mathcal{E}_2] A * [U] := \\ H * L \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_2 A * U$$

for some **sidecondition** L
residue U

Hints

$$\frac{H * [L] \models [\mathcal{E}_3 \Rightarrow \mathcal{E}_2] A * [U]}{\Delta, H \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_2 A * G}$$

Hints

$$\frac{H * [L] \Vdash [\mathcal{E}_3 \Rightarrow \mathcal{E}_2] A * [U] \quad \Delta \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_3 L * (U * G)}{\Delta, H \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_2 A * G}$$

Hints

$$\frac{H * [L] \Vdash [\mathcal{E}_3 \Rightarrow \mathcal{E}_2] A * [U] \quad \Delta \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_3 L * (U \multimap G)}{\Delta, H \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_2 A * G}$$

No backtracking: once a hint is found, we stick with it!

Hints

how to detect?

$$\frac{H * [L] \models [\mathcal{E}_3 \Rightarrow \mathcal{E}_2] A * [U] \quad \Delta \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_3 L * (U * G)}{\Delta, H \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_2 A * G}$$

Hints

$H * [L] \Vdash [\mathcal{E}_3 \Rrightarrow \mathcal{E}_2] A * [U]$ is a typeclass

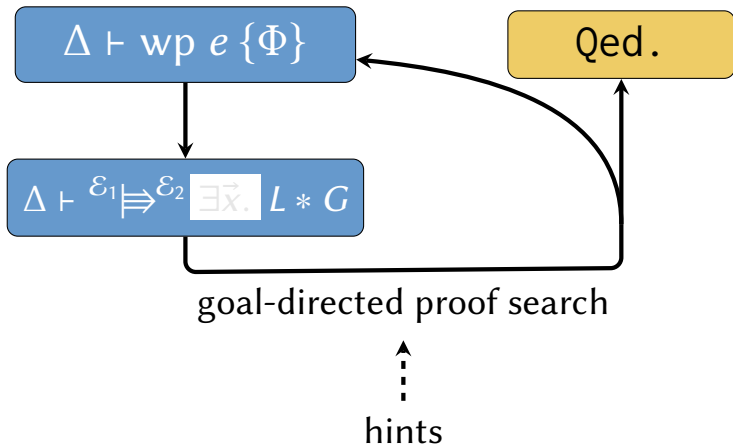
- ▶ Diaframe base hints
- ▶ Language-specific hints
- ▶ Libraries with ghost-theory hints

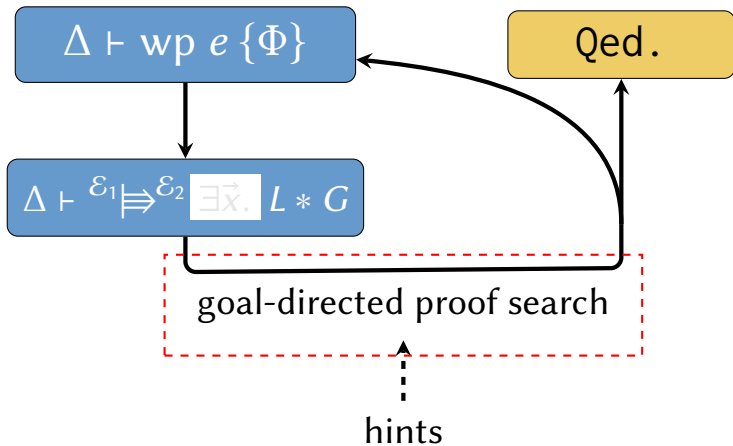
Hints

$H * [L] \Vdash [\mathcal{E}_3 \Rightarrow \mathcal{E}_2] A * [U]$ is a typeclass

Recursive rules for additional instances:

$$\frac{H * [L_2] \Vdash [\mathcal{E}_3 \Rightarrow \mathcal{E}_2] A * [U]}{(L_1 \multimap H) * [L_1 * L_2] \Vdash [\mathcal{E}_3 \Rightarrow \mathcal{E}_2] A * [U]}$$





Proof search

$\forall \vec{x}. \{L\} e \{ \Psi \}$ atomic e

$$\frac{\Delta \vdash \top \stackrel{?}{\Rightarrow} \exists \vec{x}. L * (\forall v. \Psi v \rightarrow \stackrel{?}{\Rightarrow} \top \text{wp } K[v] \{ \Phi \})}{\Delta \vdash \text{wp } K[e] \{ \Phi \}}$$

combines WP-BIND, WP-ATOMIC and WP-WAND

Proof search

$$\top \Vdash^? \exists \vec{x}. L * (\forall v. \Psi v \rightarrow^* \Vdash^{\top} \text{wp } K[v] \{\Phi\})$$

proceed like Lithium

postpone introduction when possible

Proof search

Like Lithium:

$$\frac{\forall x. \quad \Delta \vdash G}{\Delta \vdash \forall x. G}$$

$$\frac{\Delta \vdash H_1 \multimap (H_2 \multimap G)}{\Delta \vdash (H_1 * H_2) \multimap G}$$

$$\frac{\Delta \vdash \forall x. (H \multimap G)}{\Delta \vdash (\exists x. H) \multimap G}$$

$$\frac{\Delta, H \vdash G}{\Delta \vdash H \multimap G}$$

Proof search

$$L ::= \ulcorner \phi \urcorner \mid A \mid L * L$$

$$\frac{}{\Delta \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_2 \mid L * G}$$

Proof search

$$L ::= \ulcorner \phi \urcorner \mid \mathbf{A} \mid L * L$$

$$\frac{H * [L] \Vdash [\mathcal{E}_3 \Rightarrow \mathcal{E}_2] A * [U] \quad \Delta \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_3 L * (U \multimap G)}{\Delta, H \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_2 A * G}$$

Proof search

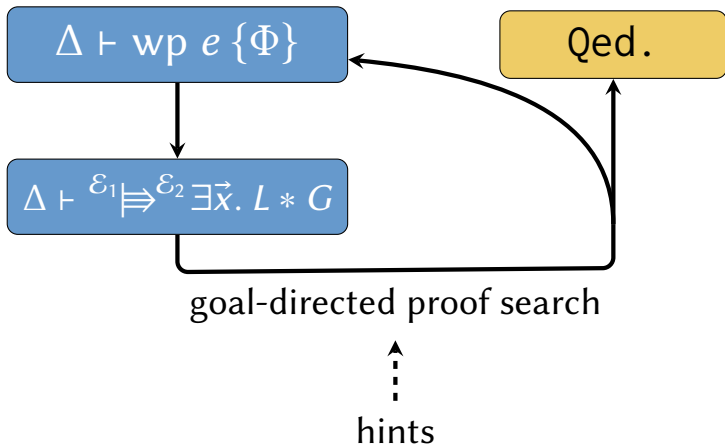
$$L ::= \ulcorner \phi \urcorner \mid A \mid L * L$$

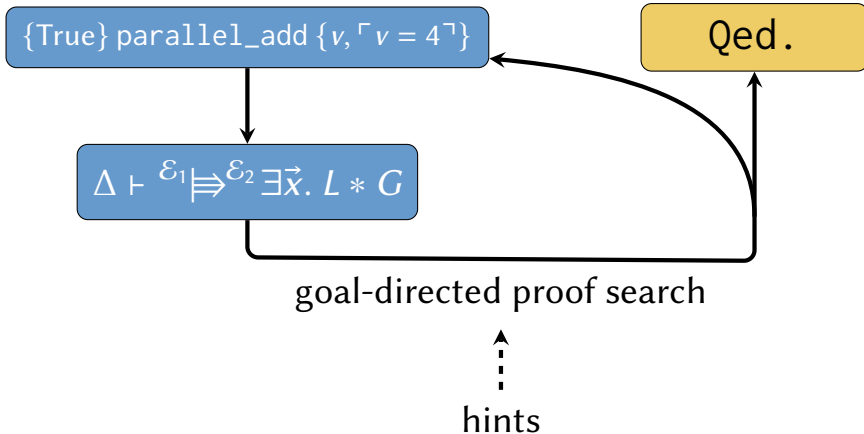
$$\frac{\Delta \vdash \varepsilon_1 \Vdash^? L_1 * \left(? \Vdash^{\varepsilon_2} L_2 * G \right)}{\Delta \vdash \varepsilon_1 \Vdash^{\varepsilon_2} (L_1 * L_2) * G}$$

Proof search

$$L ::= \boxed{\ulcorner \phi \urcorner} \mid A \mid L * L$$

$$\frac{\phi \quad \Delta \vdash G}{\Delta \vdash \varepsilon \multimap^{\varepsilon} \ulcorner \phi \urcorner * G}$$

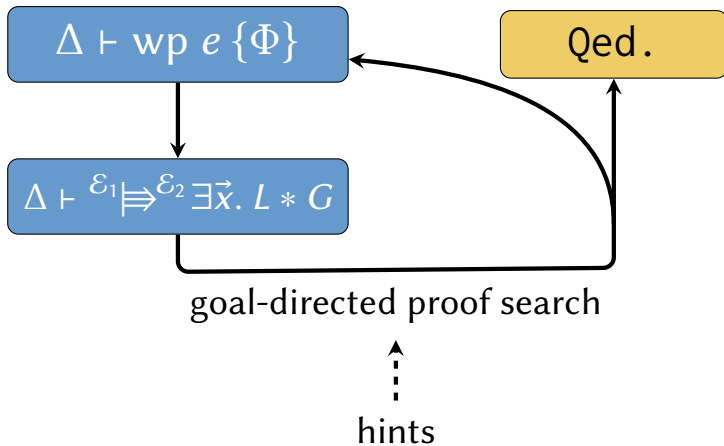




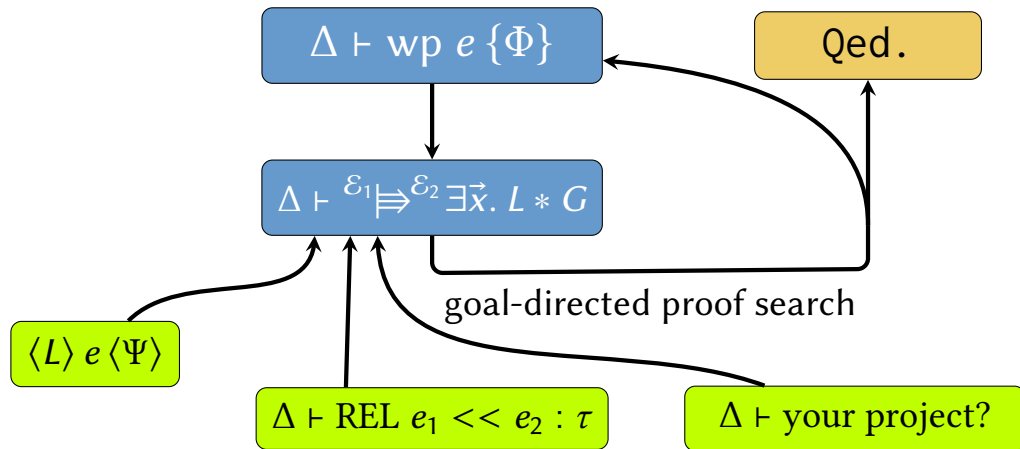
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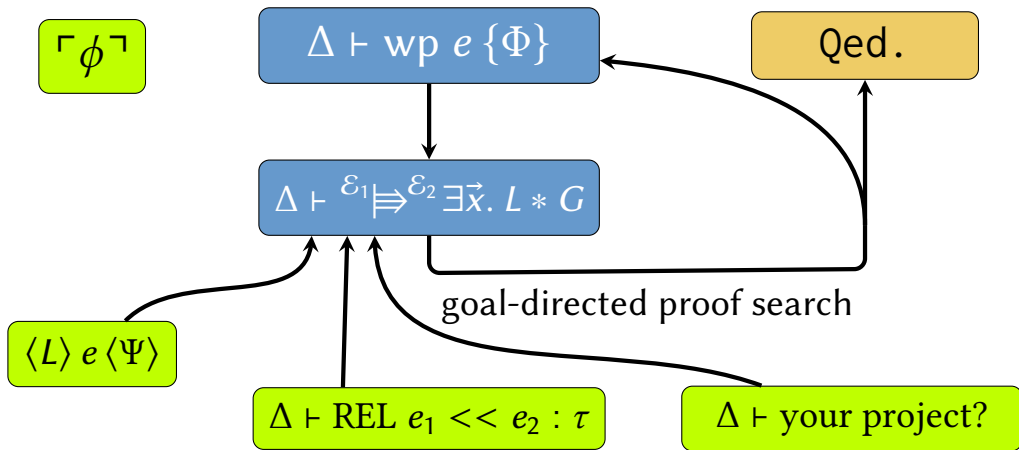
Future work



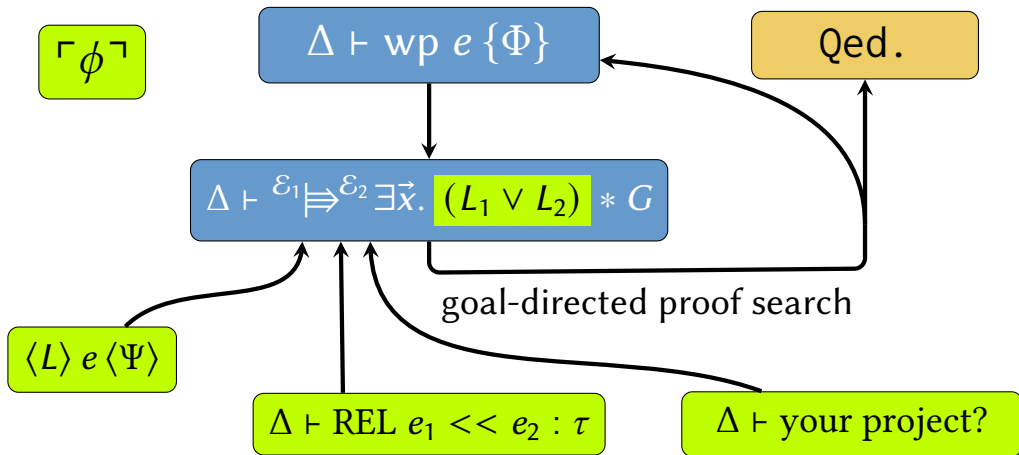
Future work



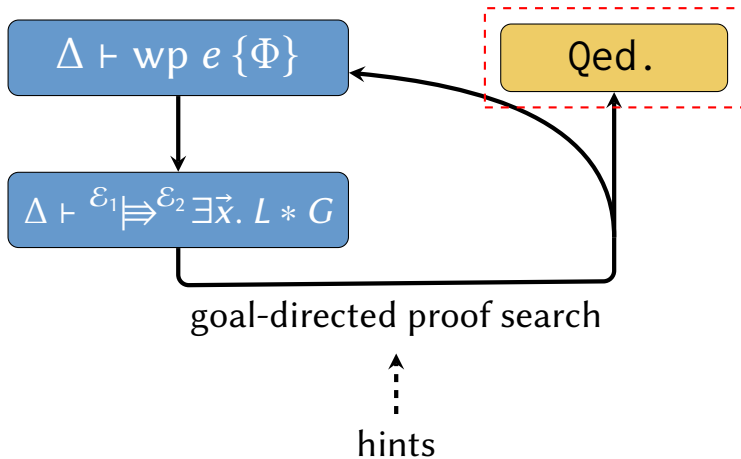
Future work



Future work



Questions?



Hint definition

$$H * [\vec{y}; L] \models [\mathcal{E}_3 \Rightarrow \mathcal{E}_2] \vec{x}; A * [U] :=$$
$$\forall \vec{y}. \quad H * L \vdash \mathcal{E}_3 \Rightarrow \mathcal{E}_2 \exists \vec{x}. A * U$$

Proof search - detailed

$$L ::= \ulcorner \phi \urcorner \mid A \mid L * L \mid \exists z. L$$

$$\frac{}{\Delta \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_2 \exists \vec{x}. L * G}$$

Proof search - detailed

$$L ::= \ulcorner \phi \urcorner \mid \mathbf{A} \mid L * L \mid \exists z. L$$

$$\frac{H * [\vec{y}; L] \models [\mathcal{E}_3 \Rightarrow \mathcal{E}_2] \vec{x}; A * [U] \quad \Delta \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_3 \exists \vec{y}. L * (\forall x. U * G)}{\Delta, H \vdash \mathcal{E}_1 \Rightarrow \mathcal{E}_2 \exists \vec{x}. A * G}$$

Proof search - detailed

$$L ::= \lceil \phi \rceil \mid A \mid L * L \mid \exists z. L$$

$$\frac{\vec{s} = \text{FV}(L_1) \quad \vec{t} = \vec{x} \setminus \vec{s} \quad \Delta \vdash \varepsilon_1 \Vdash^? \exists \vec{s}. L_1 * \left(\varepsilon_2 \Vdash^? \exists \vec{t}. L_2 * G \right)}{\Delta \vdash \varepsilon_1 \Vdash^{\varepsilon_2} \exists \vec{x}. (L_1 * L_2) * G}$$

Proof search - detailed

$$L ::= \boxed{\ulcorner \phi \urcorner} \mid A \mid L * L \mid \exists z. L$$

$$\frac{\phi[\vec{z}] \quad \Delta \vdash G[\vec{x}/\vec{z}]}{\Delta \vdash \varepsilon \rightrightarrows \varepsilon \exists \vec{x}. \ulcorner \phi \urcorner * G}$$

Proof search - detailed

$L ::= A \mid L * L \mid \exists z. L$

$$\frac{\Delta \vdash \mathcal{E}_1 \Downarrow \mathcal{E}_2 \exists(\vec{x}, t). L * G}{\Delta \vdash \exists \vec{x}. (\exists t. L) * G}$$